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# Path Planning for Point Robots

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(윤성익)

**Course URL:**  
<http://sglab.kaist.ac.kr/~sungeui/MPA>

**KAIST**



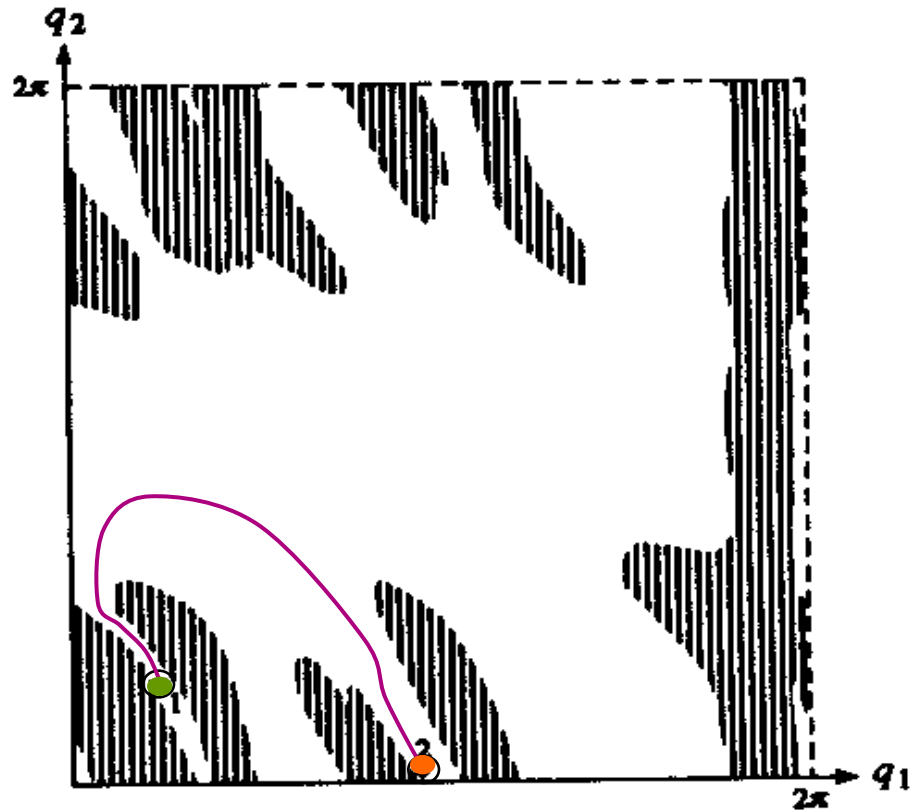
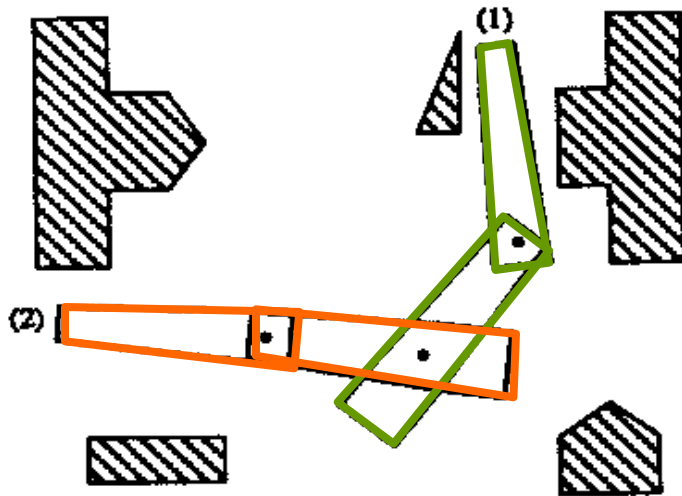
# Class Objectives

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- **Motion planning framework**
- **Classic motion planning approaches**

# Configuration Space: Tool to Map a Robot to a Point



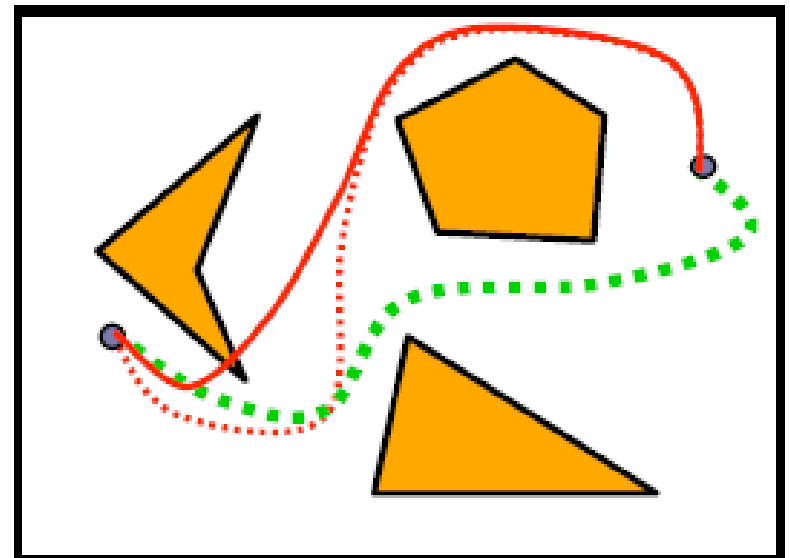
# Problem

## □ Input

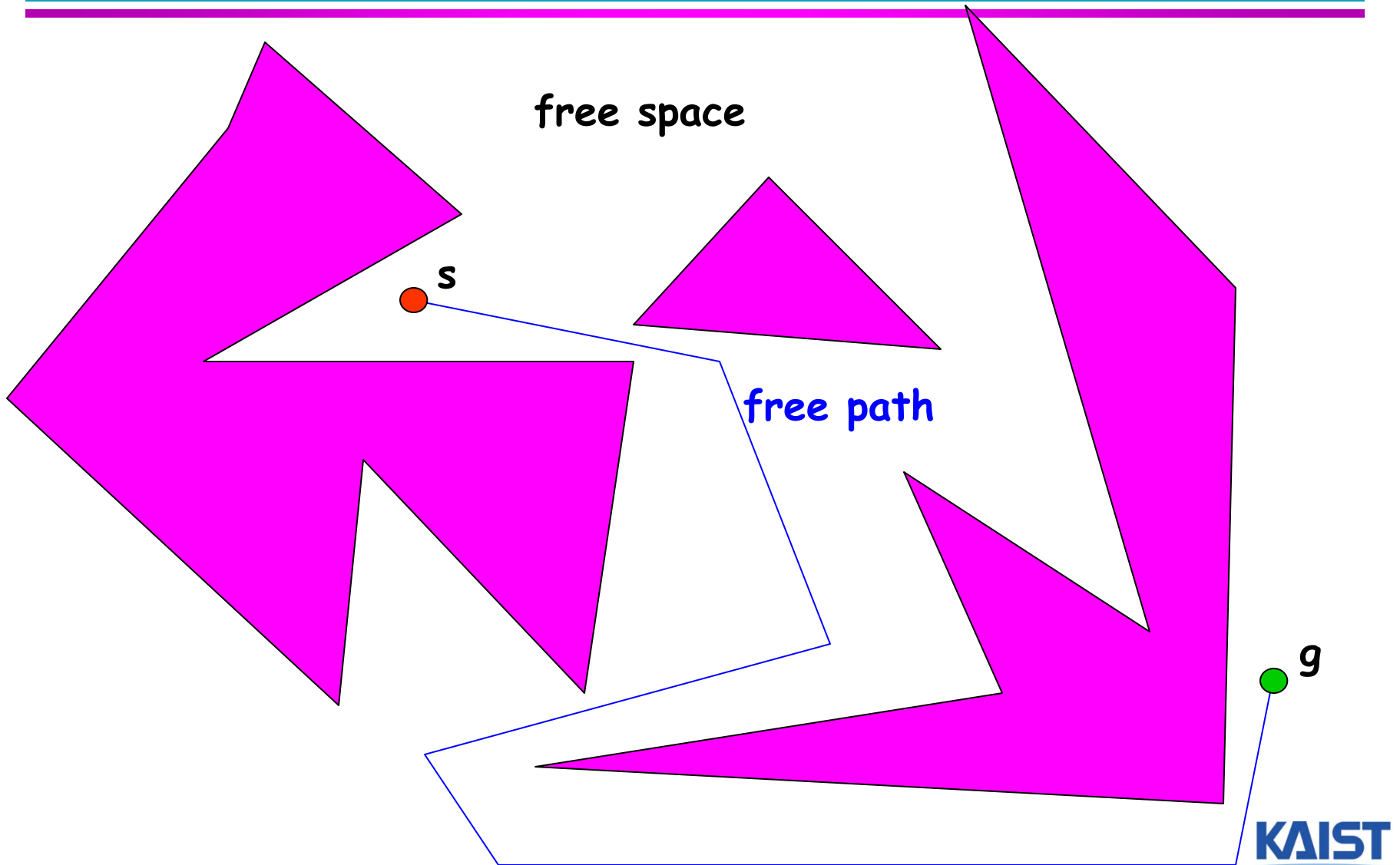
- Robot represented as a **point in the plane**
- Obstacles represented as polygons
- Initial and goal positions

## □ Output

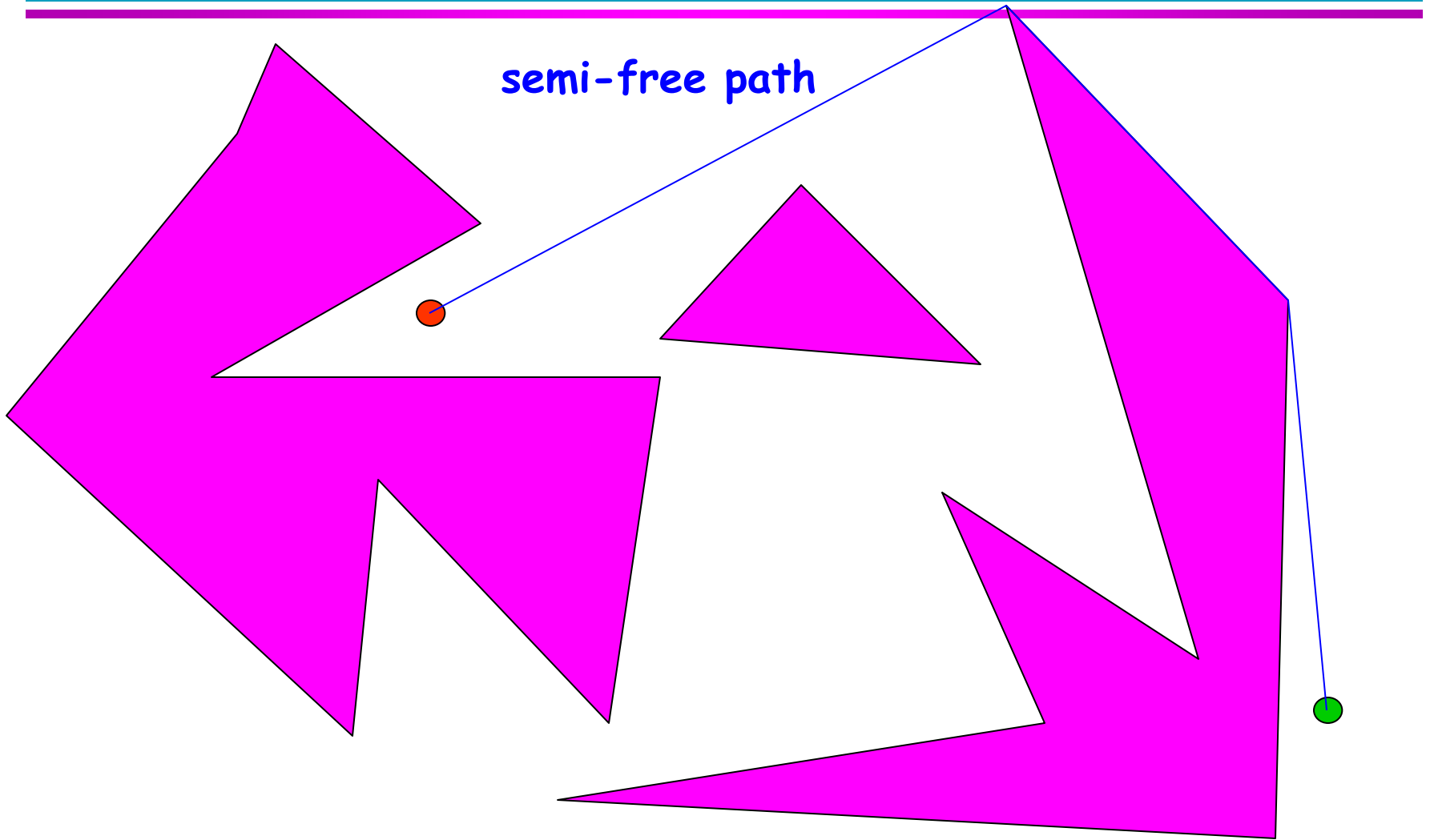
A collision-free path between the initial and goal positions



# Problem



# Problem



# Types of Path Constraints

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- **Local** constraints:  
lie in free space
- **Differential** constraints:  
have bounded curvature
- **Global** constraints:  
have minimal length

# Motion-Planning Framework

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**Continuous representation**  
(configuration space formulation)



**Discretization**  
(random sampling, processing critical geometric events)

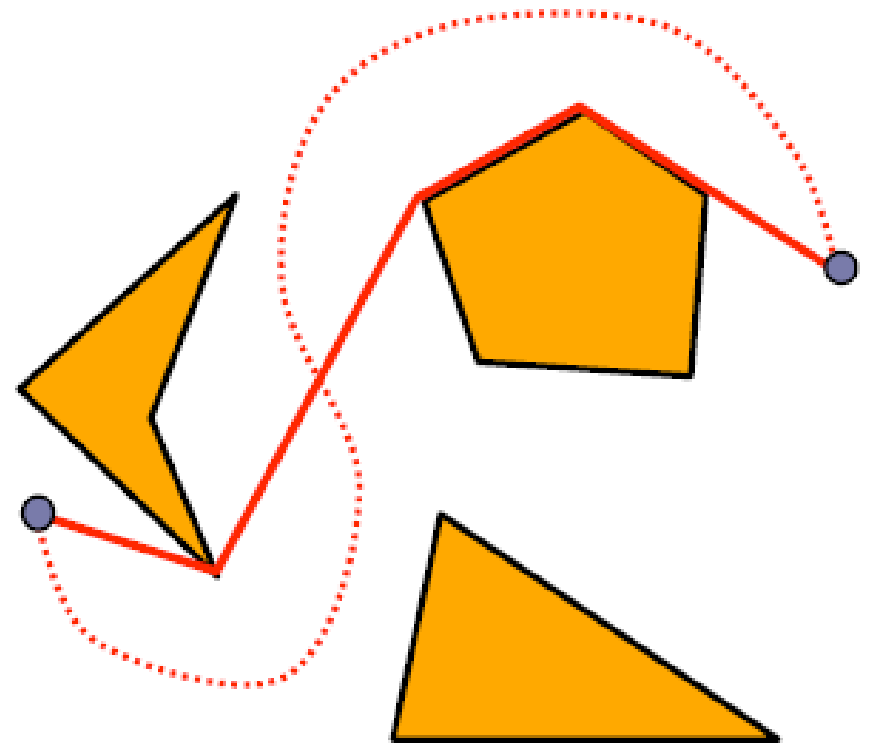


**Graph searching**  
(blind, best-first, A\*)



# Visibility graph method

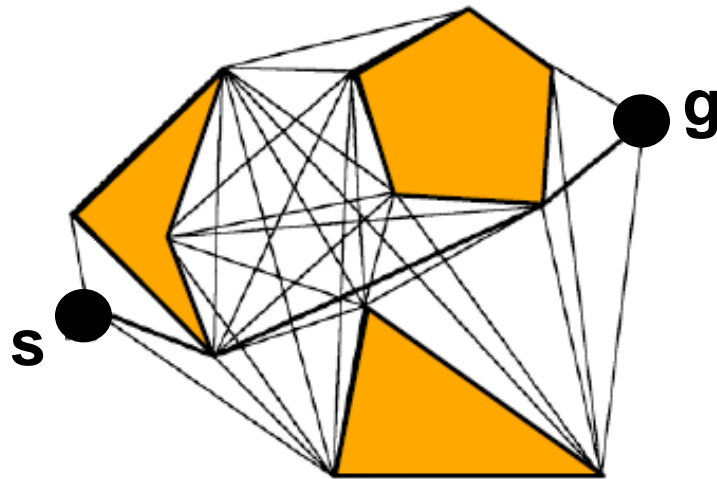
- **Observation:** If there is a collision-free path between two points, then there is a polygonal path that bends only at the obstacle vertices.
- **Why?**  
Any collision-free path can be transformed into a polygonal path that bends only at the obstacle vertices.
- A **polygonal path** is a piecewise linear curve.



# Visibility Graph

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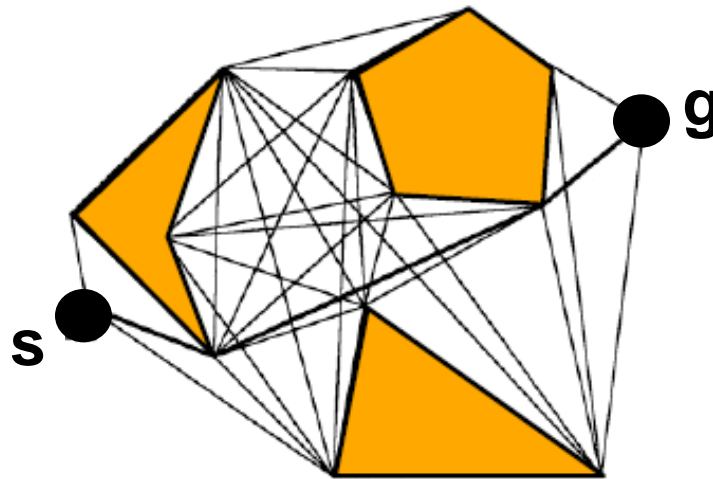


- A **visibility graph** is a graph such that
  - Nodes:  $s$ ,  $g$ , or obstacle vertices
  - Edges: An edge exists between nodes  $u$  and  $v$  if the line segment between  $u$  and  $v$  is an obstacle edge or it does not intersect the obstacles

# Visibility Graph

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- A **visibility graph**
  - Introduced in the late 60s
  - Can produce shortest paths in 2-D configuration spaces

# Simple Algorithm

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- **Input:**  $s, q$ , polygonal obstacles
  - **Output:** visibility graph  $G$
- 1: **for** every pair of nodes  $u, v$
  - 2:   **if** segment  $(u, v)$  is an obstacle **then**
  - 3:     insert edge  $(u, v)$  into  $G$ ;
  - 4:   **else**
  - 5:     **for** every obstacle edge  $e$
  - 6:       **if** segment  $(u, v)$  intersects  $e$
  - 7:         go to (1);
  - 8:     insert edge  $(u, v)$  into  $G$ ;
  - 9: Search a path with  $G$  using  $A^*$

# Computation Efficiency

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1: <b>for</b> every pair of nodes $u, v$	$O(n^2)$
2: <b>if</b> segment $(u, v)$ is an obstacle <b>then</b>	$O(n)$
3:     insert edge $(u, v)$ into $G$ ;	
4: <b>else</b>	
5: <b>for</b> every obstacle edge $e$	$O(n)$
6: <b>if</b> segment $(u, v)$ intersects $e$	
7:         go to (1);	
8:     insert edge $(u, v)$ into $G$ ;	

- **Simple algorithm:  $O(n^3)$  time**
- **More efficient algorithms**
  - Rotational sweep  $O(n^2 \log n)$  time, etc.
- **$O(n^2)$  space**

# Motion-Planning Framework

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**Continuous representation**  
(configuration space formulation)



**Discretization**  
(random sampling, processing critical geometric events)



**Graph searching**  
(blind, best-first, A\*)

# Graph Search Algorithms

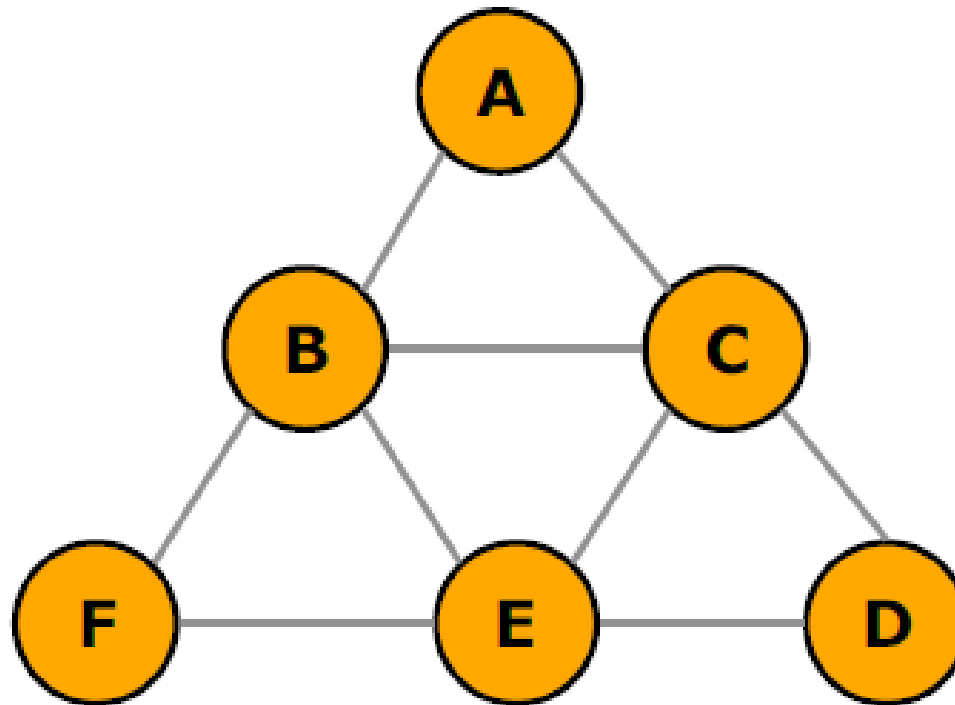
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- Breadth, depth-first, best-first
- Dijkstra's algorithm
- A\*

# Breadth-first search

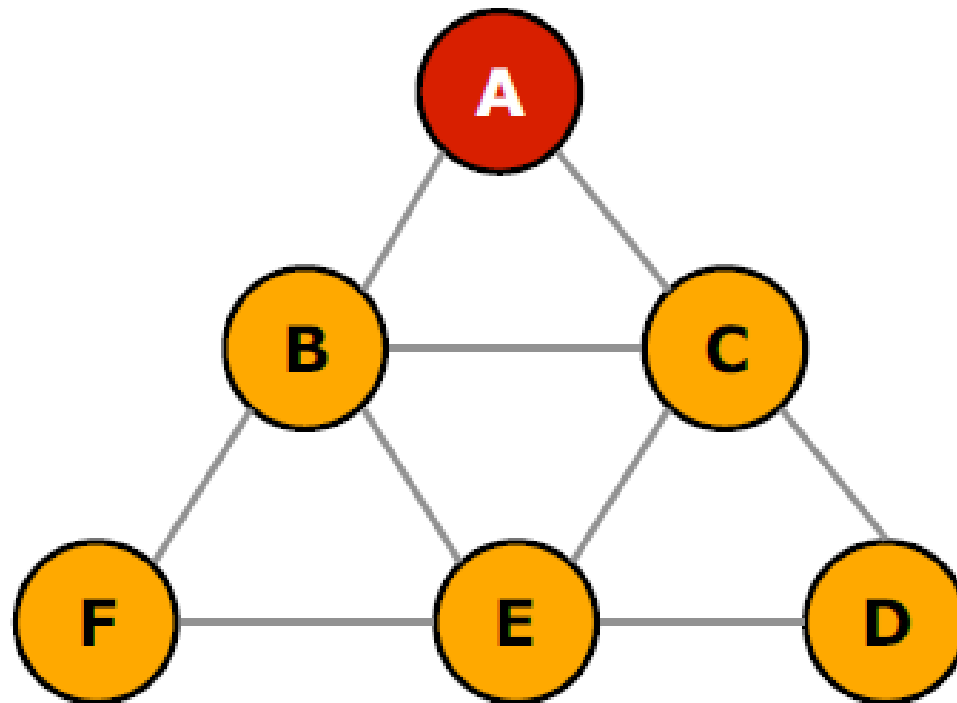
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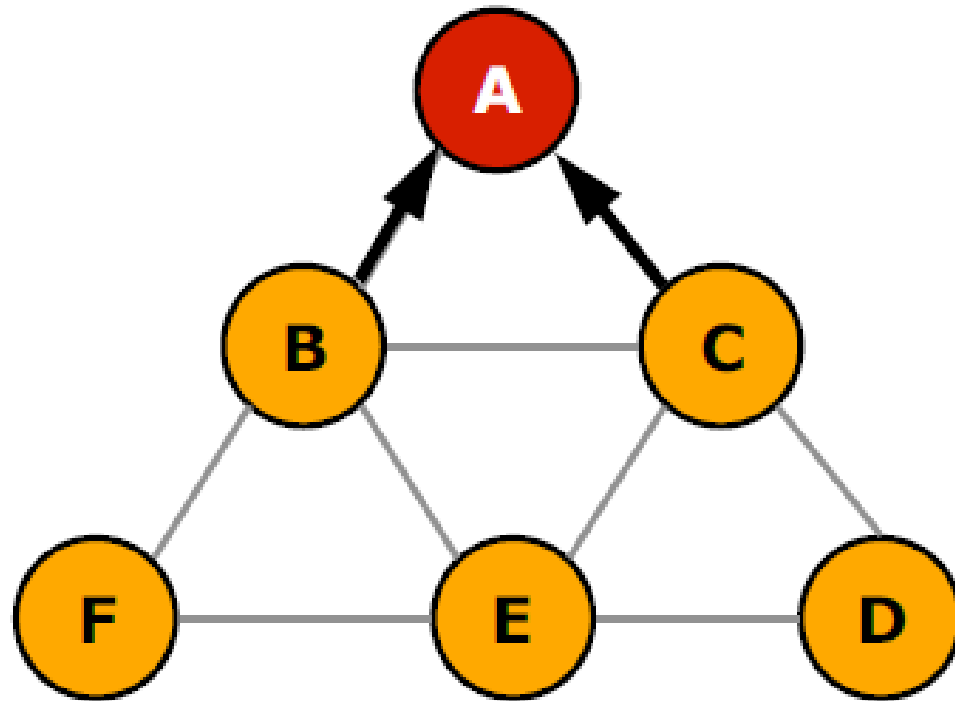
# Breadth-first search

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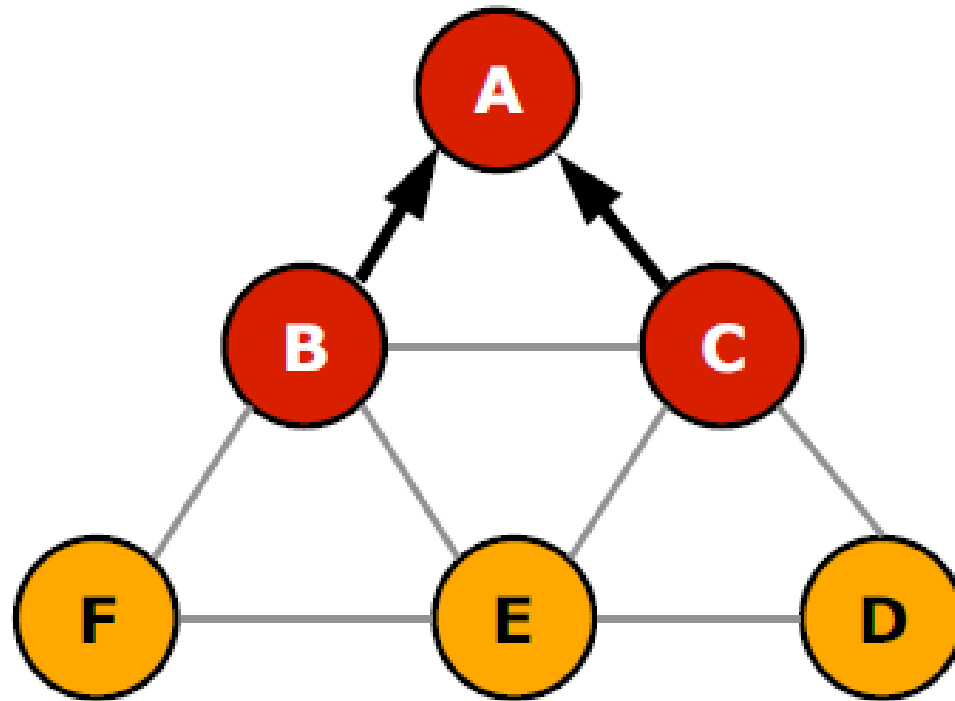
# Breadth-first search

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# Breadth-first search

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# Dijkstra's Shortest Path Algorithm

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- Given a (non-negative) weighted graph, two vertices,  $s$  and  $g$ :
  - Find a path of minimum total weight between them
  - Also, find minimum paths to other vertices
  - Has  $O(|V| \lg|V| + |E|)$

# Dijkstra's Shortest Path Algorithm

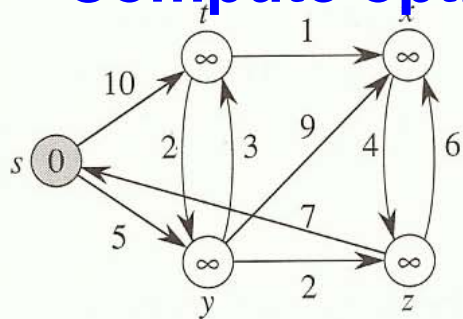
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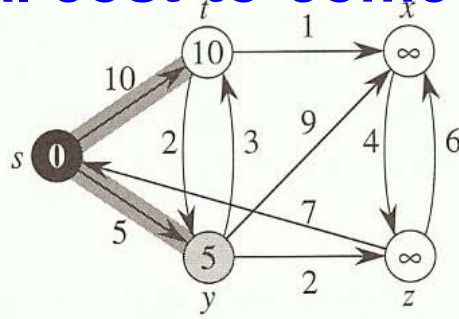
- A set  $S$ 
  - Contains vertices whose final shortest-path cost has been determined
- **DIJKSTRA** ( $G, s$ )
  1. Initialize-Single-Source ( $G, s$ )
  2.  $S \leftarrow \text{empty}$
  3. Queue  $\leftarrow$  Vertices of  $G$
  4. **While** Queue is not empty
  5.     **Do**  $u \leftarrow \text{min-cost from Queue}$
  6.          $S \leftarrow \text{union of } S \text{ and } \{u\}$
  7.         **for** each vertex  $v$  in Adj [ $u$ ]
  8.             **do** RELAX ( $u, v$ )

# Dijkstra's Shortest Path Algorithm

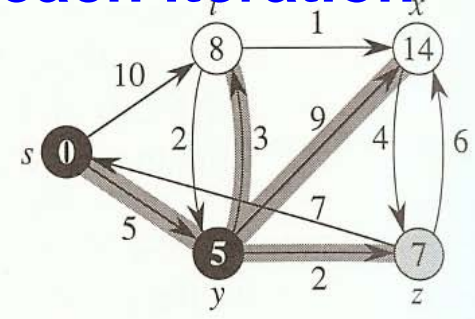
Compute optimal cost-to-come at each iteration



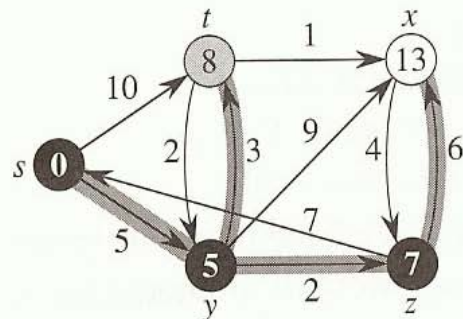
(a)



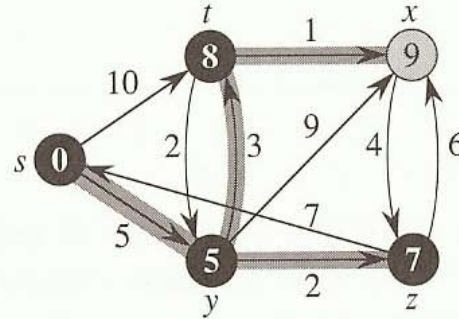
(b)



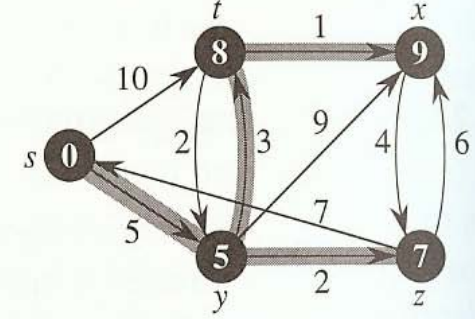
(c)



(d)



(e)



(f)

Black vertices are in the set.

White vertices are in the queue.

Shaded one is chosen for relaxation.



# Best-First Search

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- **Pick a next node based on an estimate of the optimal cost-to-go cost**
  - Greedily finds solutions that look good
  - Solutions may not be optimal
  - Can find solutions quite fast, but can be also very slow



# Framework

continuous representation



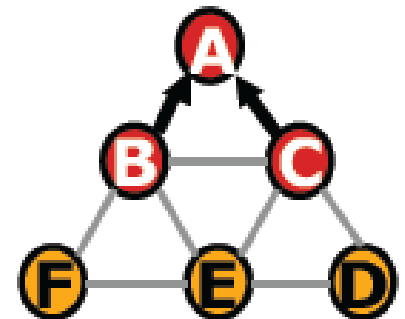
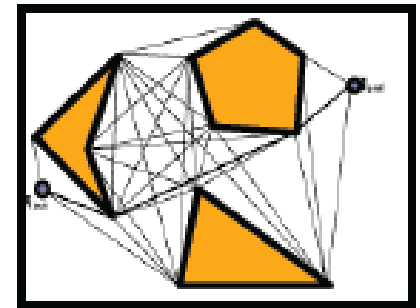
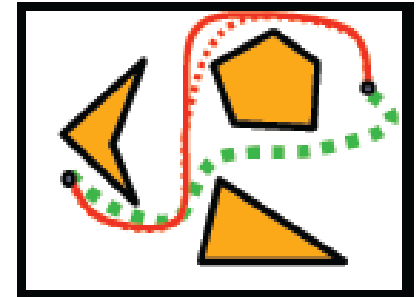
discretization

construct visibility graph



graph searching

breadth-first search



# Computational Efficiency

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- Running time  $O(n^3)$ 
  - Compute the visibility graph
  - Search the graph
- Space  $O(n^2)$
  
- Can we do better?

# Classic Path Planning Approaches

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- **Roadmap**
  - Represent the connectivity of the free space by a network of 1-D curves
- **Cell decomposition**
  - Decompose the free space into simple cells and represent the connectivity of the free space by the adjacency graph of these cells
- **Potential field**
  - Define a function over the free space that has a global minimum at the goal configuration and follow its steepest descent

# Classic Path Planning Approaches

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- **Roadmap**

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- **Cell decomposition**

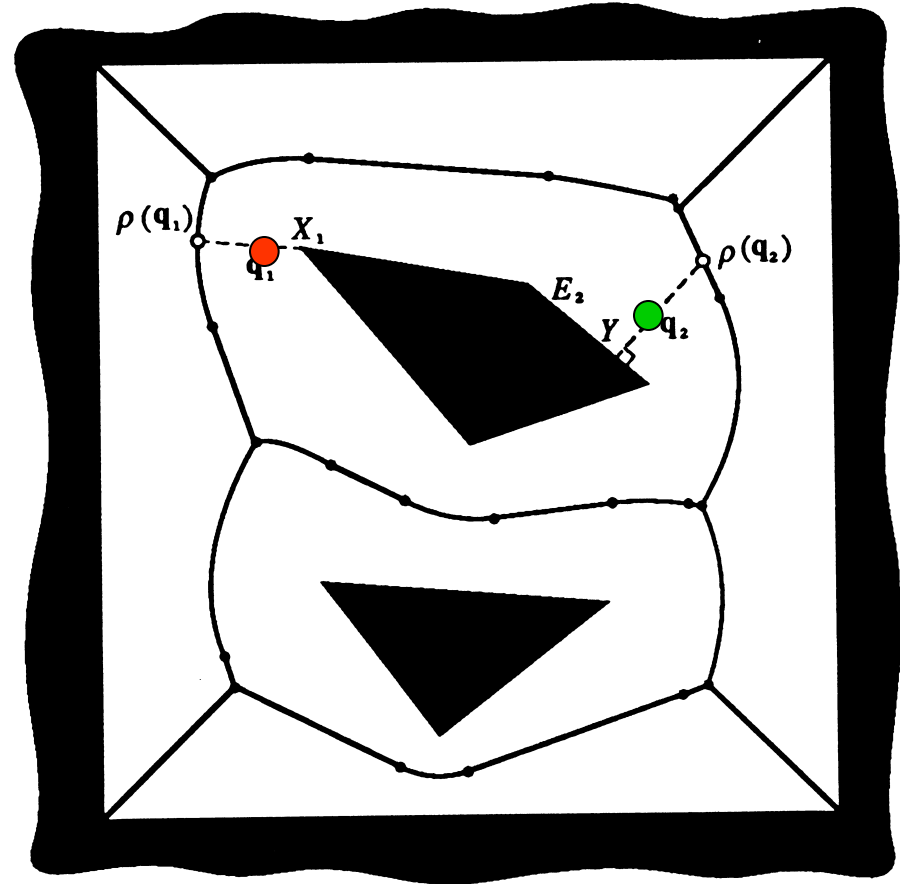
- Decompose the free space into simple cells and represent the connectivity of the free space by the adjacency graph of these cells

- **Potential field**

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# Roadmap Methods

- Visibility Graph
  - Shakey project, SRI [Nilsson 69]
- Voronoi diagram
  - Introduced by computational geometry researchers
  - Generate paths that maximize clearance
  - $O(n \log n)$  time and  $O(n)$  space



# Other Roadmap Methods

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- **Visibility graph**
- **Voronoi diagram**
- **Silhouette**
  - **First complete general method that applies to spaces of any dimension and is singly exponential in # of dimensions [Canny, 87]**
- **Probabilistic roadmaps**

# Classic Path Planning Approaches

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- **Roadmap**
  - Represent the connectivity of the free space by a network of 1-D curves
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# Cell-Decomposition Methods

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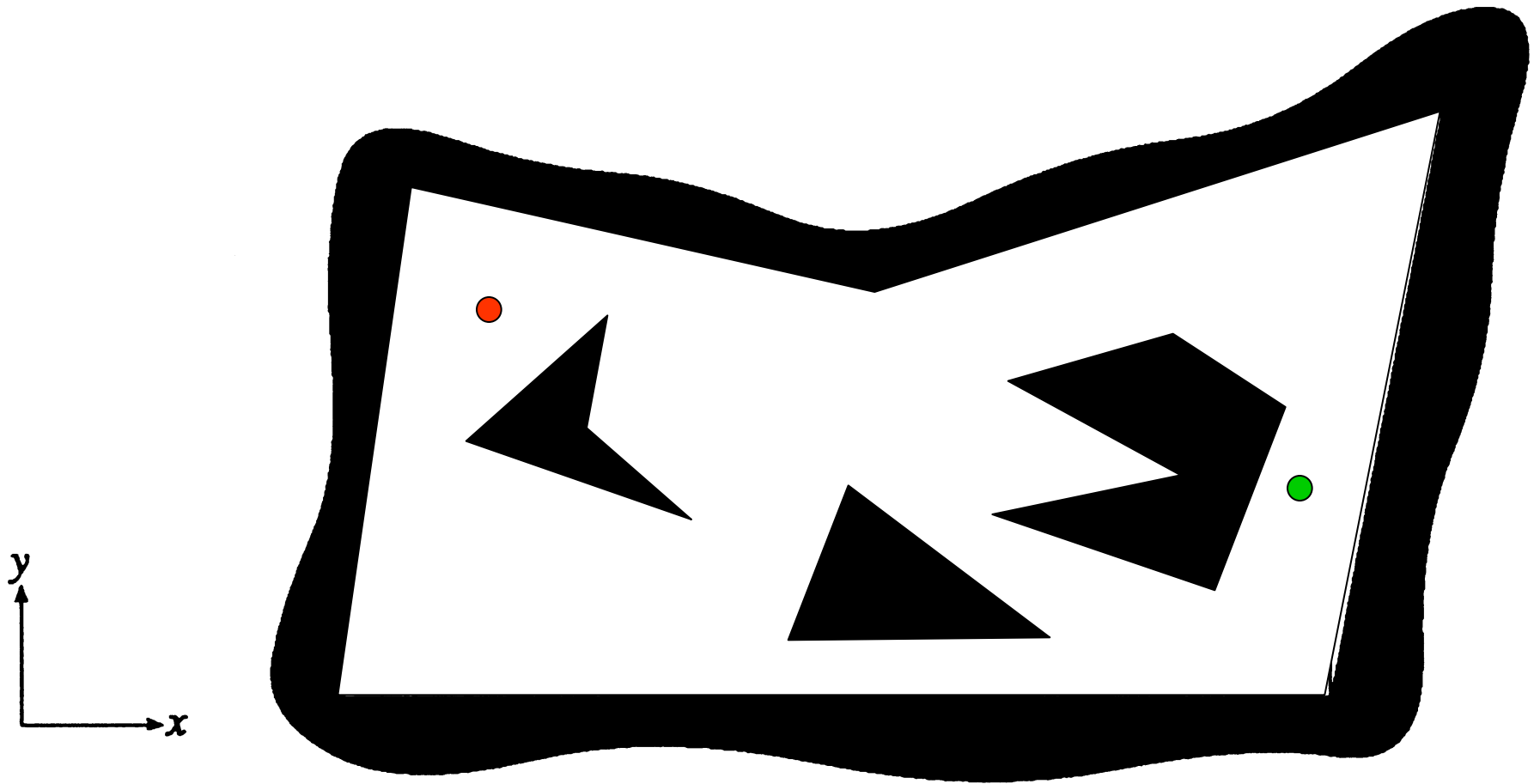
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- **Two classes of methods:**
  - **Exact and approximate cell decompositions**
  
- **Exact cell decomposition**
  - **The free space  $F$  is represented by a collection of non-overlapping cells whose union is exactly  $F$**
  - **Example: trapezoidal decomposition**



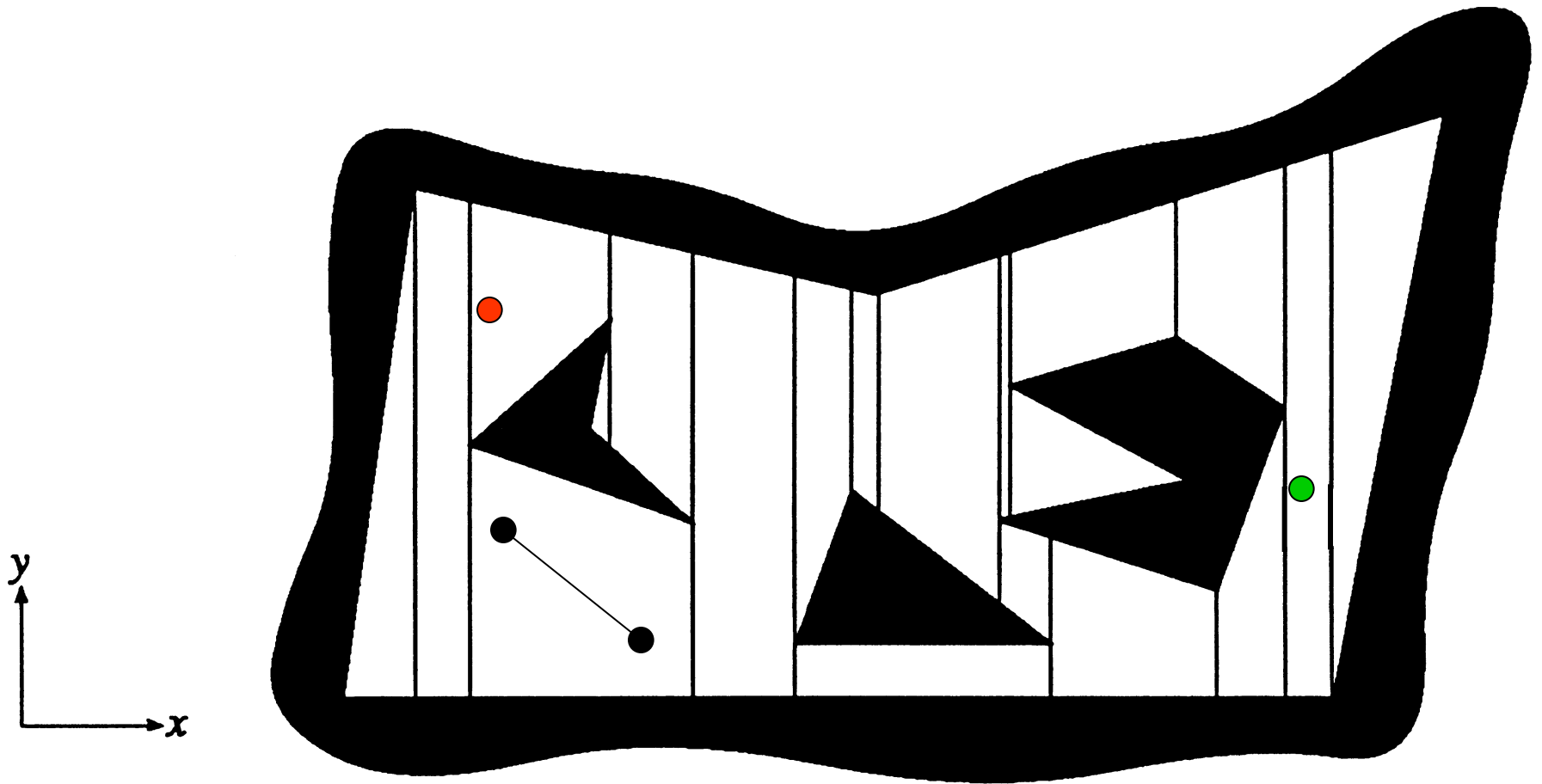
# Trapezoidal Decomposition

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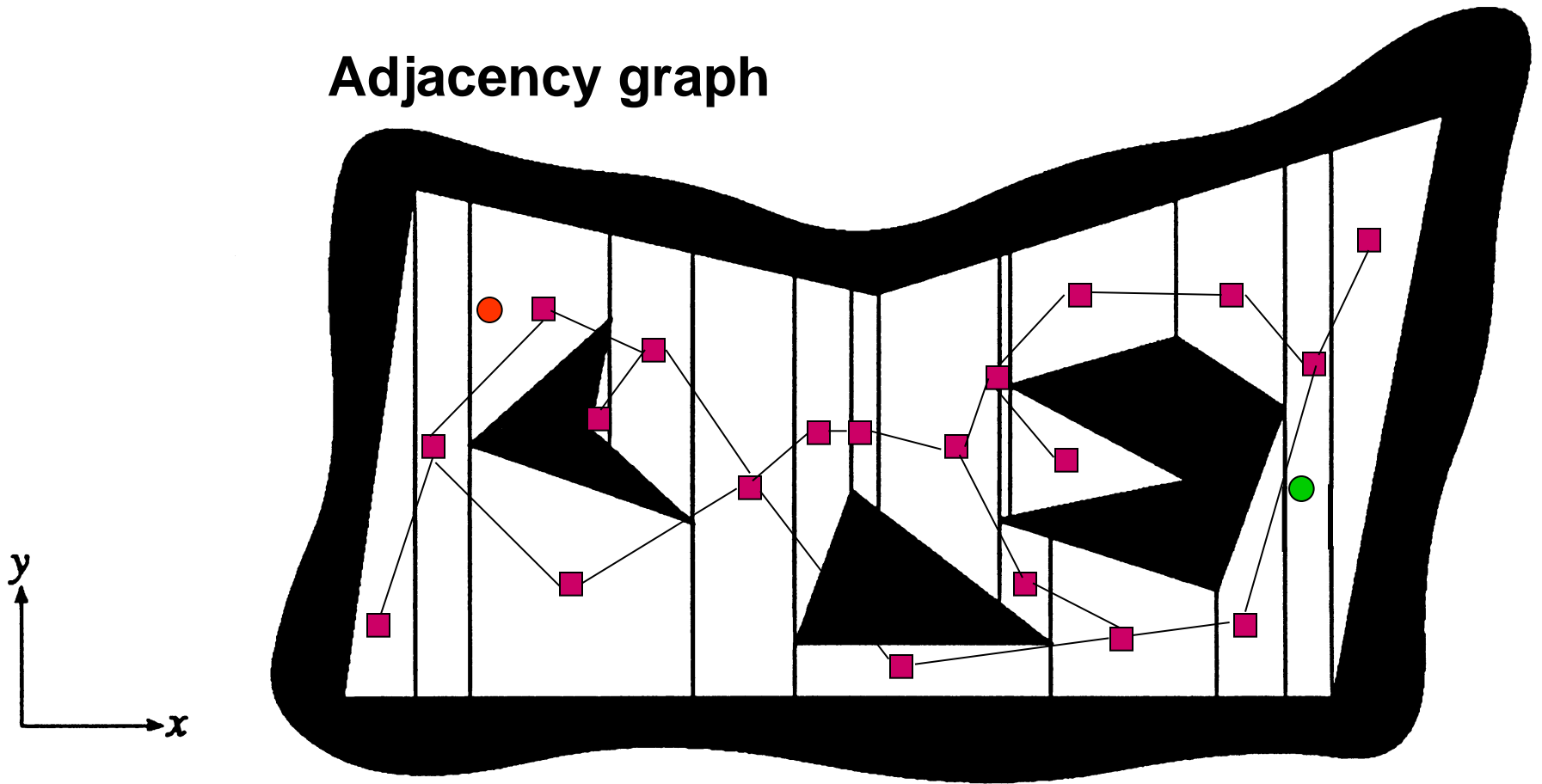
# Trapezoidal Decomposition

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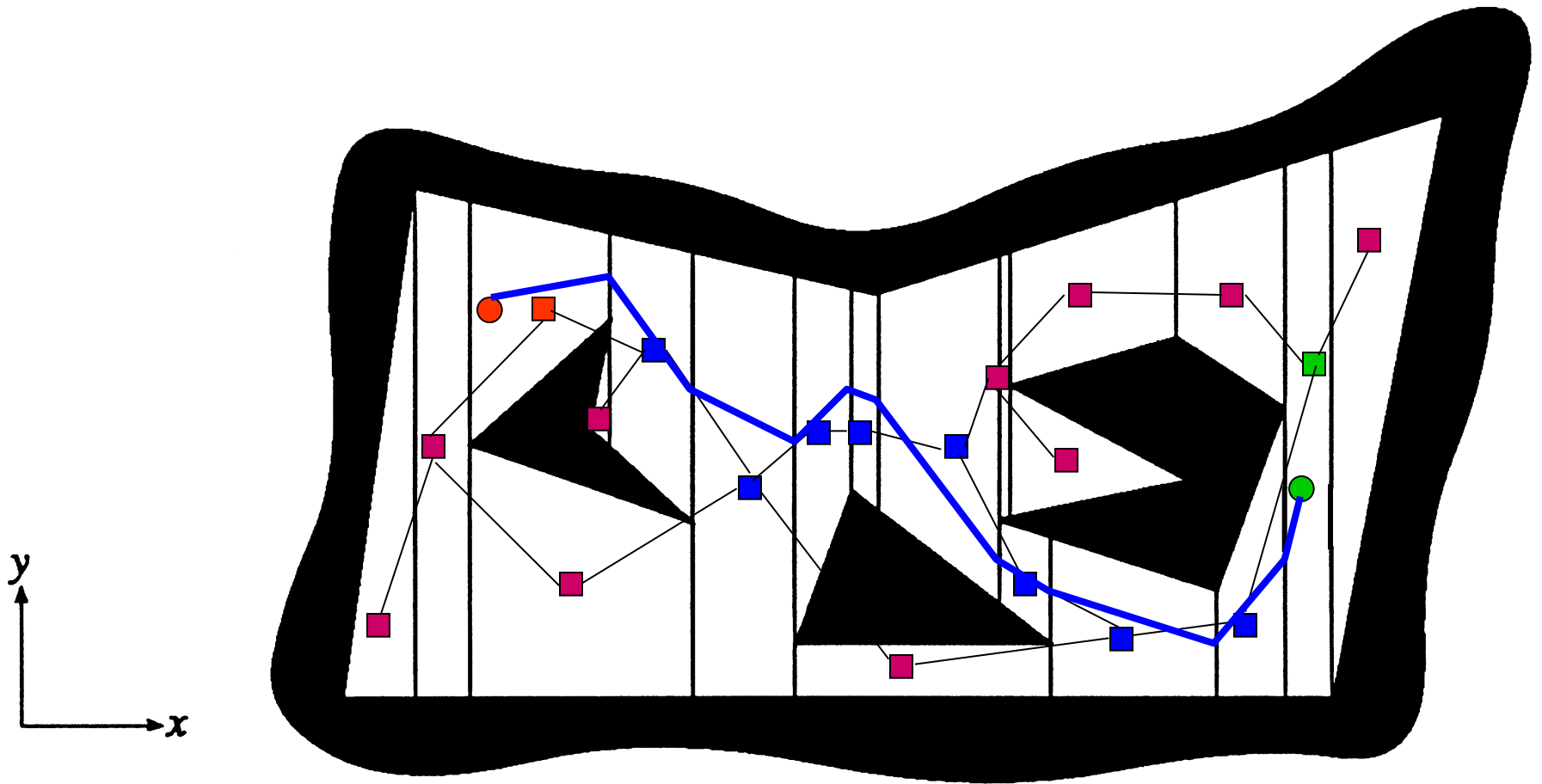


# Trapezoidal Decomposition

Adjacency graph

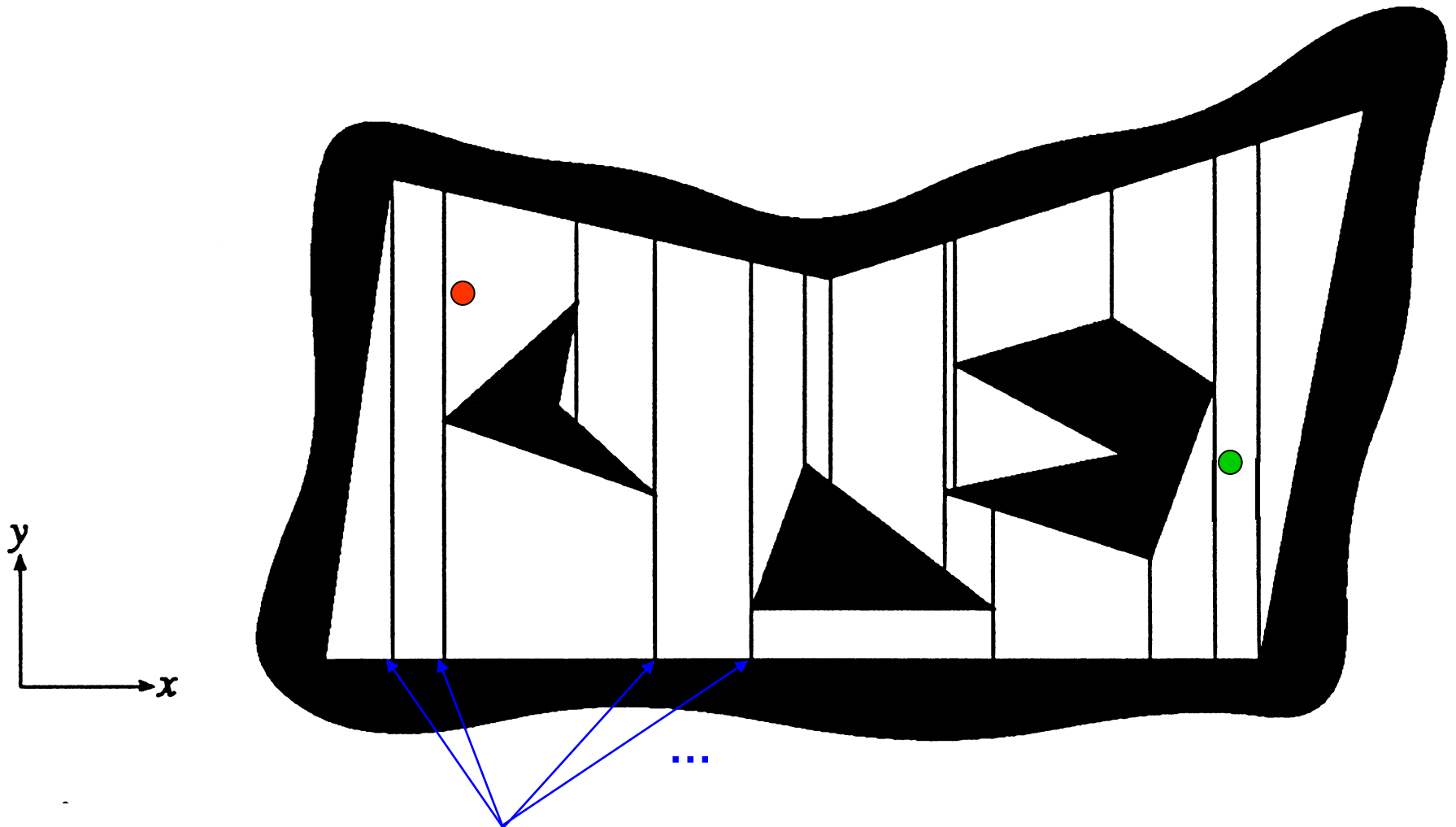


# Trapezoidal Decomposition



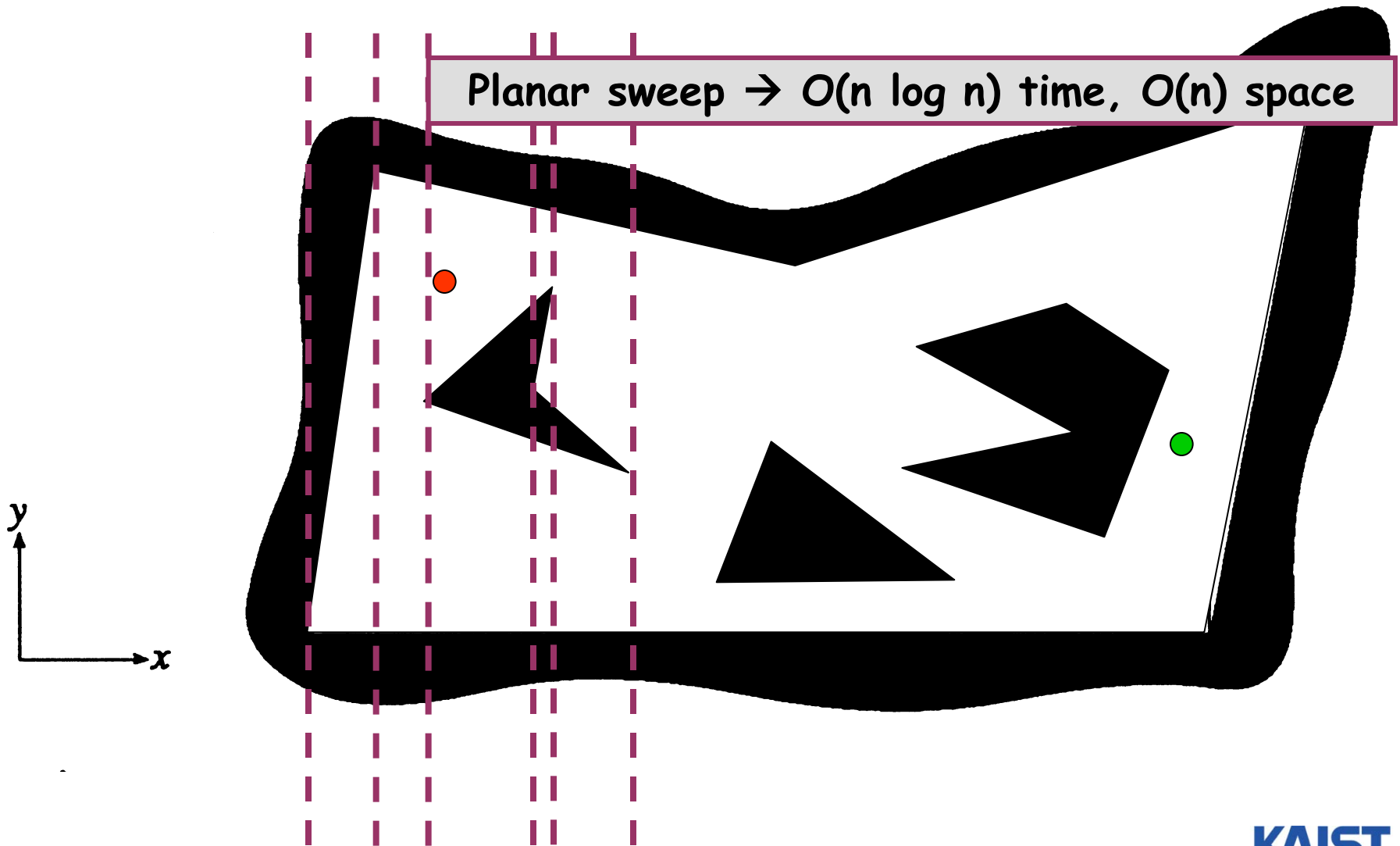
# Trapezoidal Decomposition

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critical events  $\rightarrow$  criticality-based decomposition KAIST

# Trapezoidal Decomposition



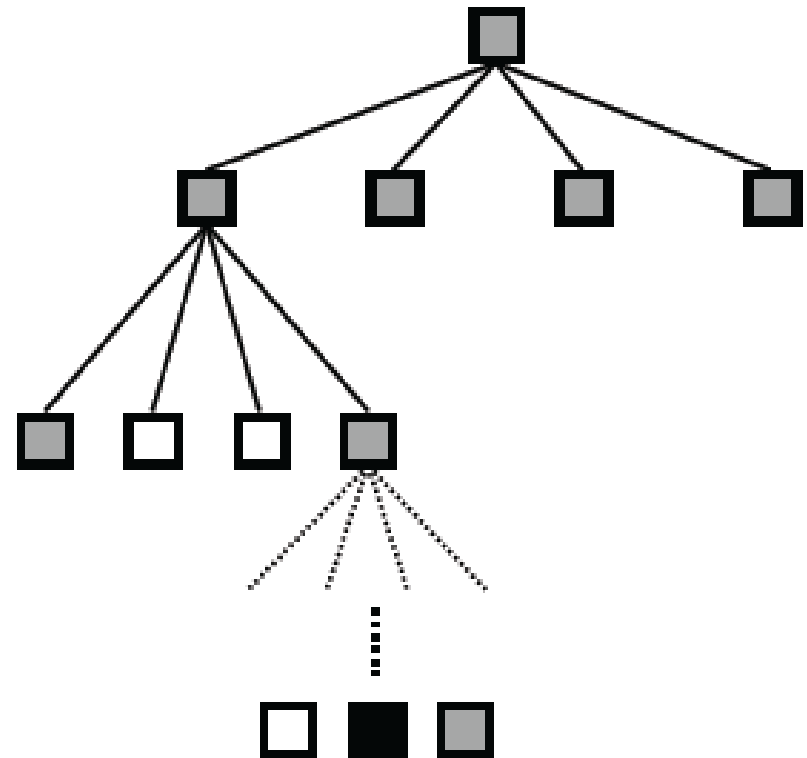
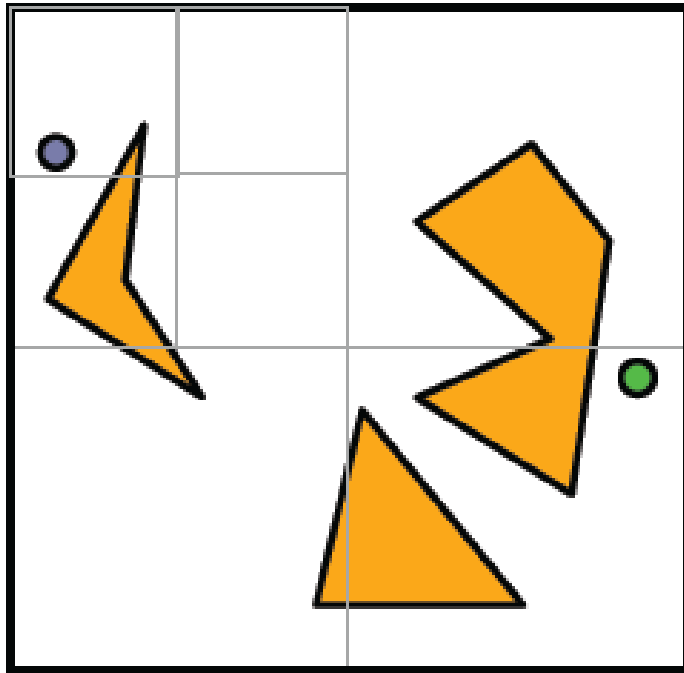
# Cell-Decomposition Methods

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- **Two classes of methods:**
  - **Exact and approximate cell decompositions**
- **Approximate cell decomposition**
  - **The free space  $F$  is represented by a collection of non-overlapping cells whose union is contained in  $F$**
  - **Cells usually have simple, regular shapes (e.g., rectangles and squares)**
  - **Facilitates hierarchical space decomposition**

# Quadtree decomposition



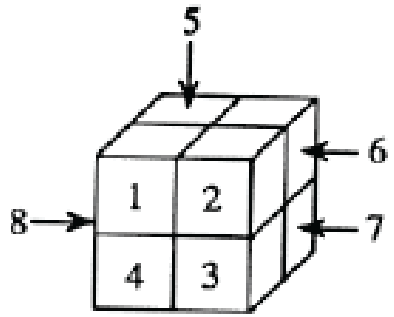
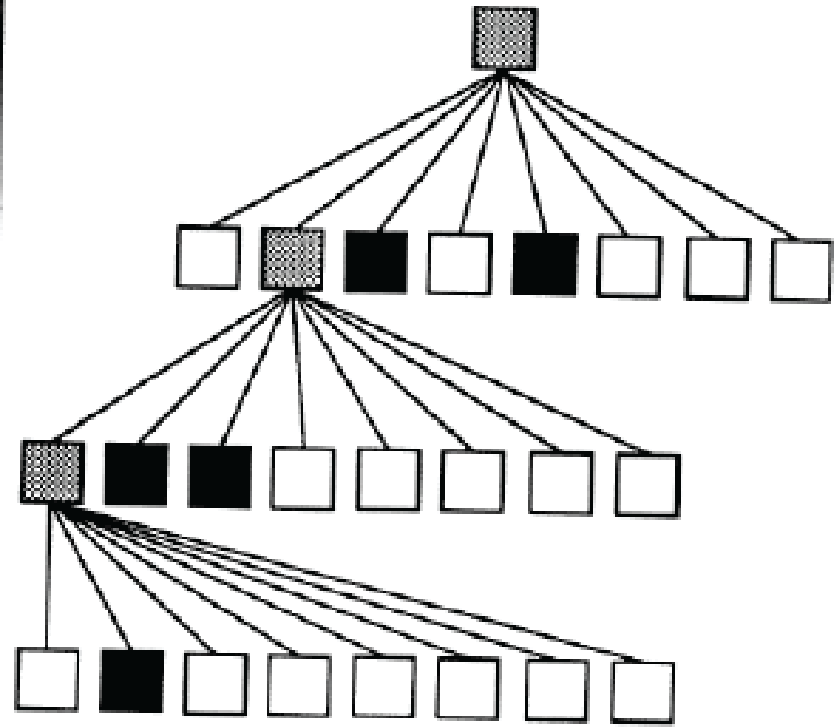
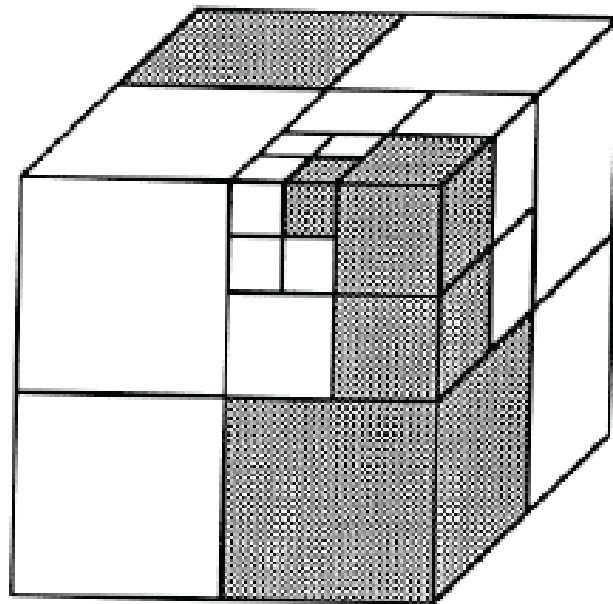
 empty

 mixed

 full



# Octree decomposition



 EMPTY cell     MIXED cell     FULL cell

# Sketch of Algorithm

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1. Decompose the free space  $F$  into cells
2. Search for a sequence of **mixed** or **free** cells that connect that initial and goal positions
3. Further decompose the mixed
4. Repeat 2 and 3 until a sequence of **free** cells is found

# Classic Path Planning Approaches

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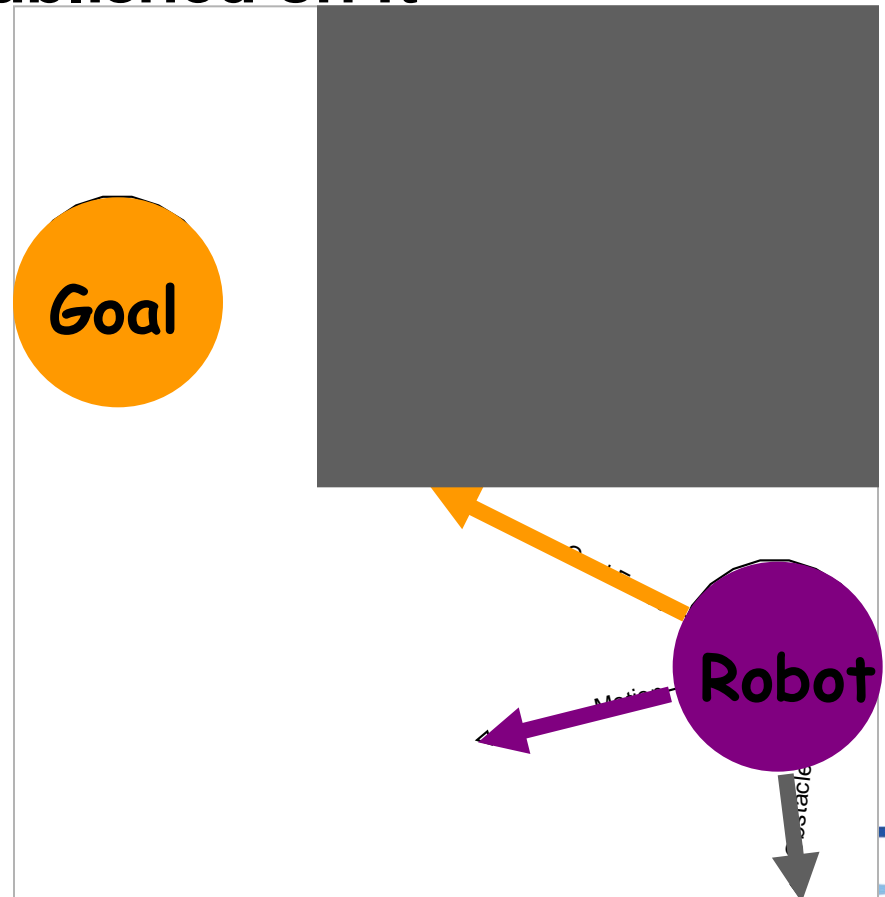
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# Potential Field Methods

- Initially proposed for real-time collision avoidance [Khatib, 86]
  - Hundreds of papers published on it

$$F_{Goal} = -k_p (x - x_{Goal})$$

$$F_{Obstacle} = \begin{cases} \eta \left( \frac{1}{\rho} - \frac{1}{\rho_0} \right) \frac{1}{\rho^2} \frac{\partial \rho}{\partial x} & \text{if } \rho \leq \rho_0, \\ 0 & \text{if } \rho > \rho_0 \end{cases}$$



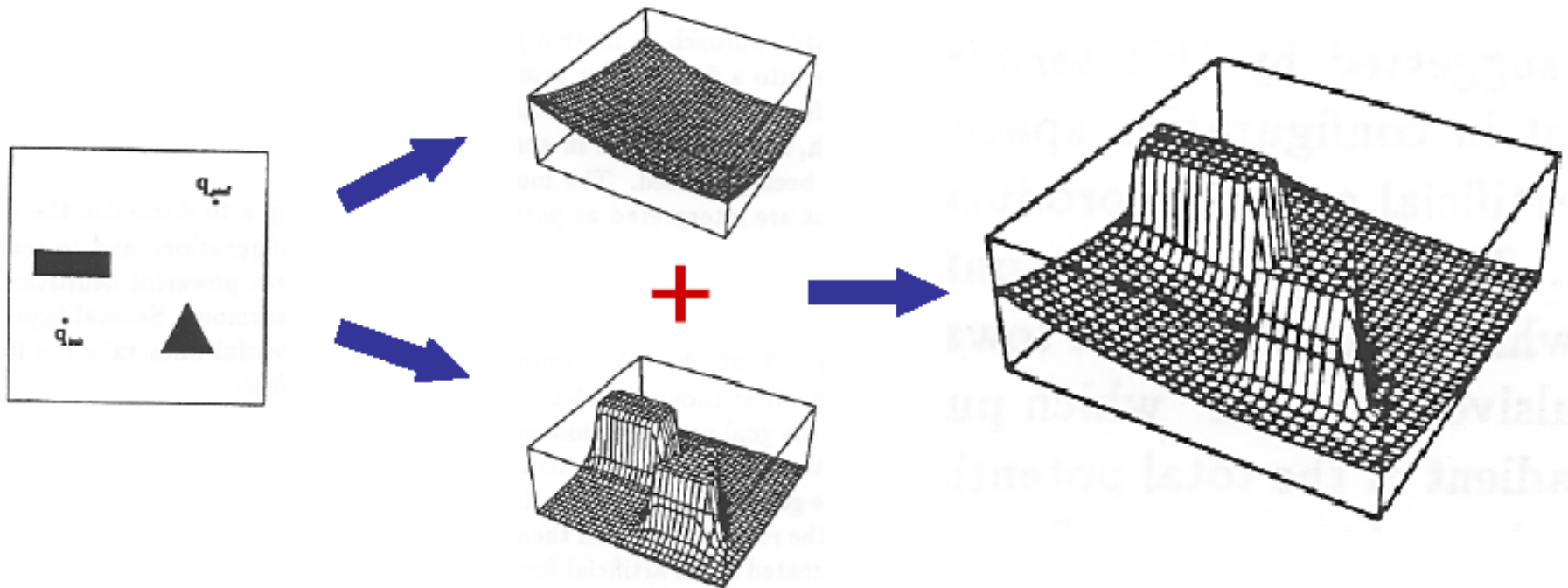
# Potential Field

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- A scalar function over the free space
- To navigate the robot applies a force proportional to the negated gradient of the potential field
- A navigation function is an ideal potential field that
  - Has global minimum at the goal
  - Has no local minima
  - Grows to infinity near obstacles
  - Is smooth

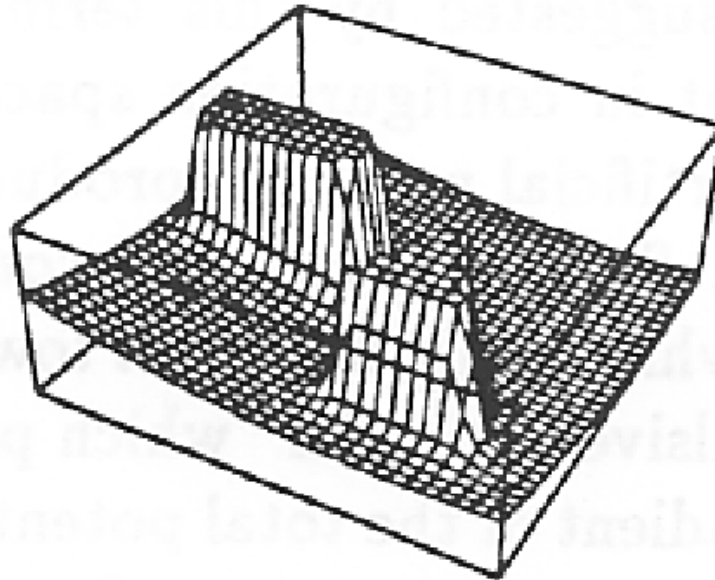
# Attractive and Repulsive fields



# Local Minima

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- **What can we do?**
  - **Escape from local minima by taking random walks**
  - **Build an ideal potential field that does not have local minima**

# Sketch of Algorithm

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- Place a regular grid  $G$  over the configuration space
- Compute the potential field over  $G$
- Search  $G$  using a best-first algorithm with potential field as the heuristic function



# Question

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- **Can such an ideal potential field be constructed efficiently in general?**

# Completeness

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- A **complete** motion planner always returns a solution when one exists and indicates that no such solution exists otherwise
  - Is the visibility algorithm complete? Yes
  - How about the exact cell decomposition algorithm and the potential field algorithm?

# Class Objectives were:

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- Motion planning framework
- Classic motion planning approaches

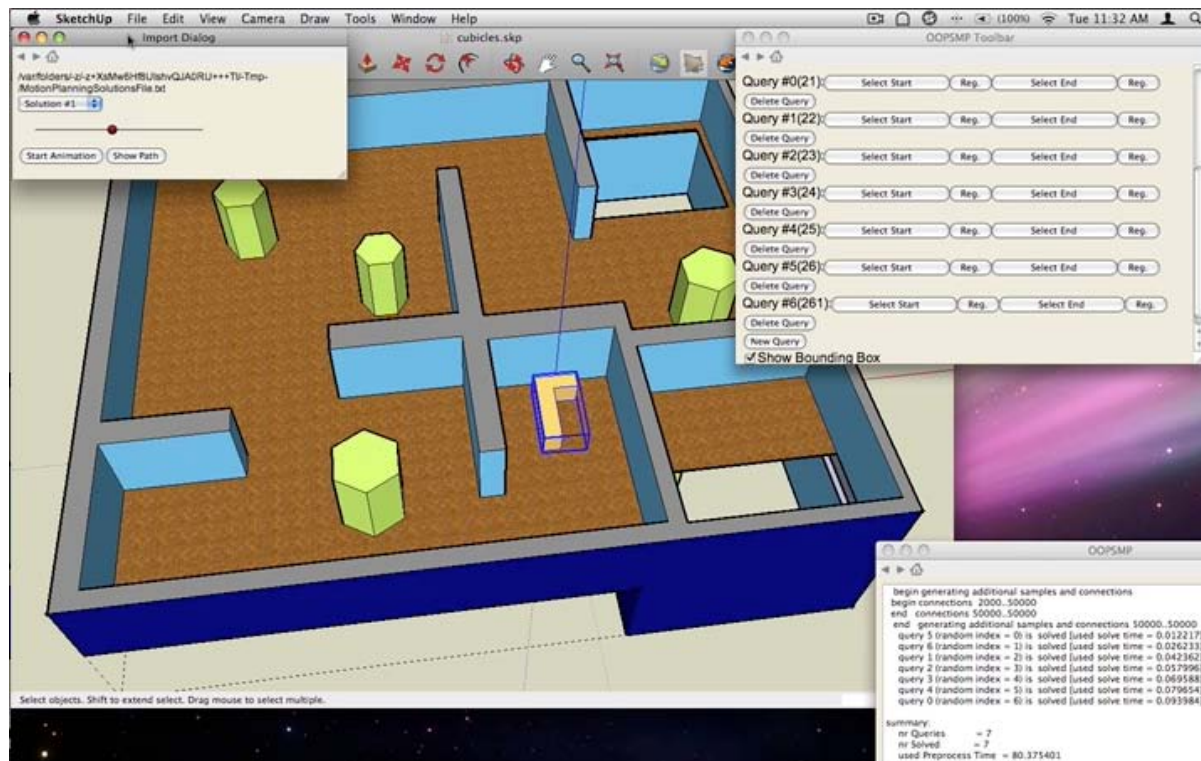
# Homework for Every Class

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- **Go over the next lecture slides**
- **Come up with one question on what we have discussed today and submit at the beginning of the next class**
  - **1**
  - **2 for typical questions**
  - **3 for questions with thoughts**
  - **4 for questions that surprised me**

# Homework

- Install [OOPSMP Motion Planning Library](#)
- Create a scene and a robot
- Find a collision-free path and visualize the path



# Conf. Deadline

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- ICRA
  - Sept. 15, 2010
- IROS
  - Feb. 28, 2010

**IROS  
2010**

take place in Taiwan



## Welcome to IROS 2010

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2010 IEEE/RSJ International Conference on Intelligent Robots and Systems

*Taipei International Convention Center, Taipei, Taiwan*

**October 18-22, 2010**

### Important Deadlines:

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*February 28, 2010*

*Proposals for Organized Sessions / Proposals for Tutorials/Workshops /  
Submission of full-length papers and videos*

# Next Time....

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- Configuration spaces